Semantics for Rollback-Based Continuous/Discrete Simulation

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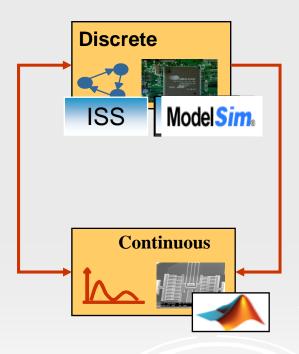
Continuous/Discrete Systems

- Applications to various domains
 - Defense, medical, communications, automotive, ...
- Examples of Continuous/Discrete Systems
 - MEMS, real-time controllers, mixed-signal systems ...
- Main characteristics
 - Complexity, heterogeneity
- Main design challenges
 - Global specification and validation

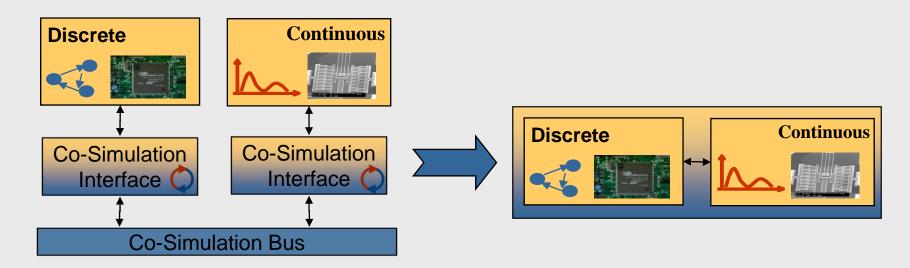


Continuous/Discrete Systems Design

- Collaboration between different teams
- Incremental refinements through different abstraction levels with specific execution models
- Validation requires joint execution of heterogeneous execution models
 - Co-Simulation Technique



Challenges for Continuous/Discrete Co-Simulation



- Defining new tools facilitating cooperation between different teams
 - Enabling easy specification, automatic generation for simulation interfaces
 - Taking into account implementation choices
 - Exploiting powerful existing tools (Simulink, SystemC, ...)
 - Based on a single well defined formalism for domain interaction

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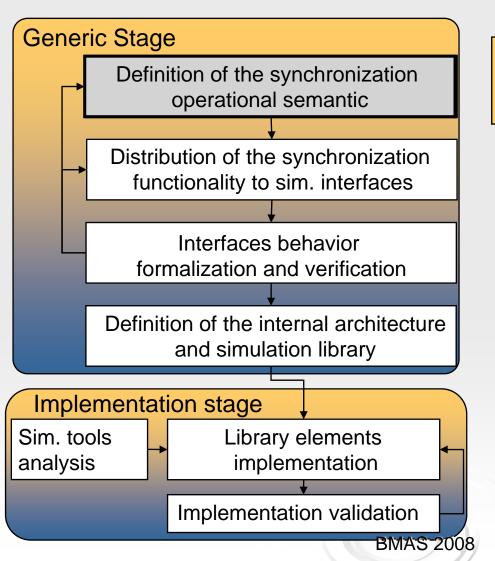
Contributions

- Definition of the semantics for the continuous and the discrete co-simulation interfaces
 - The interfaces representation using DEVS models [University of Arizona] and timed automata
- Formal verification of the simulation interfaces
 - Study for rollback-based continuous/discrete simulation models

Outline

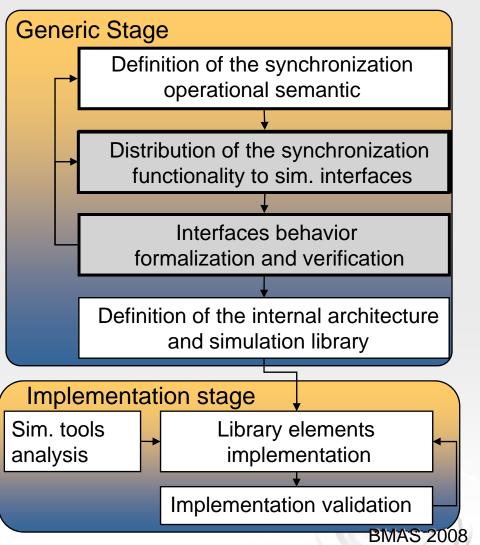
- Global Simulation for Continuous/Discrete Systems
- Design Methodology for Continuous/Discrete Systems Co-Simulation Tools
- Semantics for Rollback-based Continuous/Discrete Co-Simulation
- Formal verification for Co-Simulation Interfaces
- Conclusion

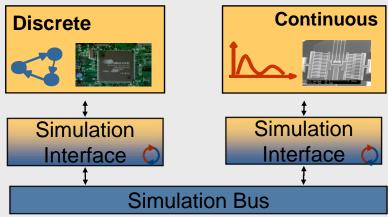
Design Methodology for Continuous/Discrete Simulation Tools



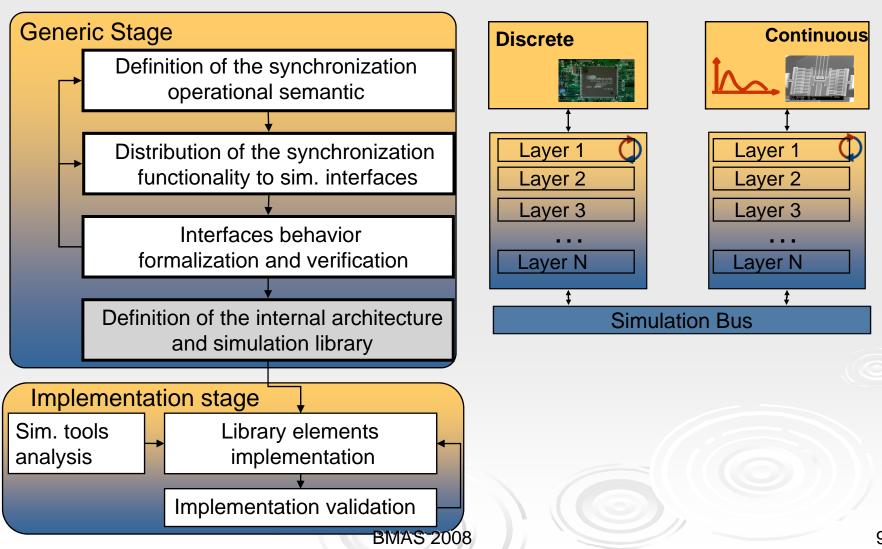


Design Methodology for Continuous/Discrete Simulation Tools

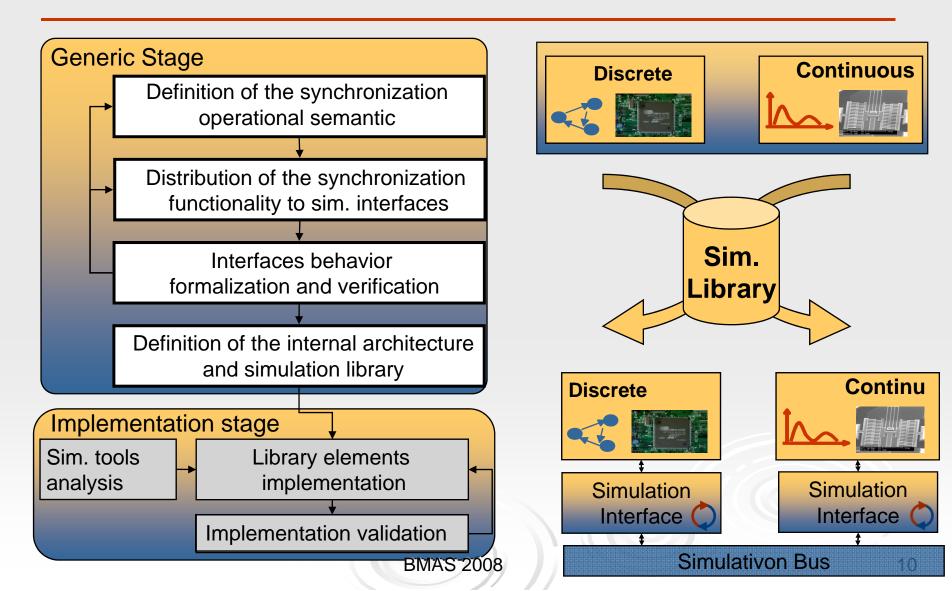




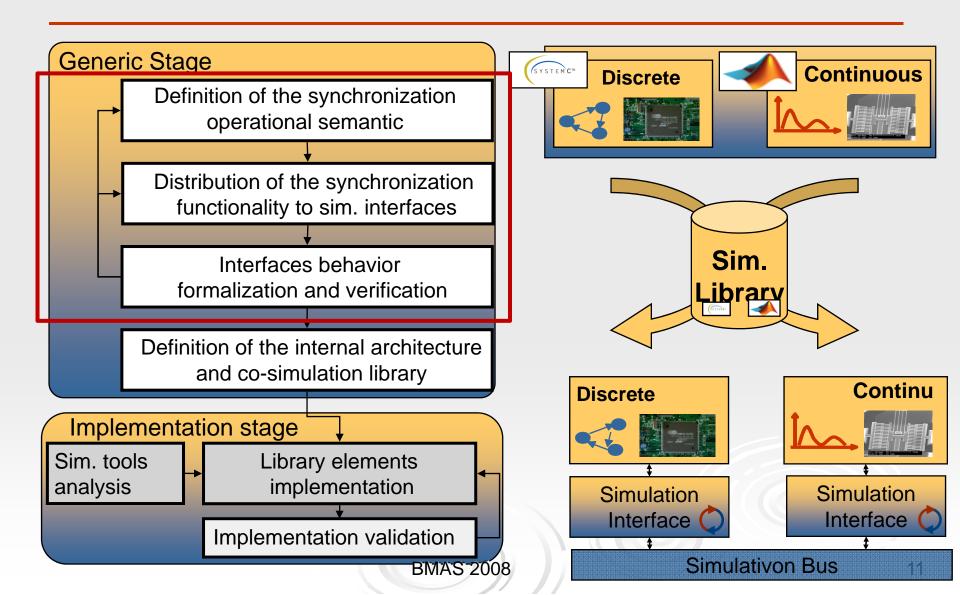
Design Methodology for Continuous/Discrete Simulation Tools



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Design Methodology for Continuous/Discrete Simulation Tools



Outline

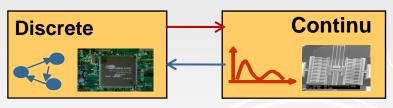
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Continuous/Discrete Systems - Global Simulation -

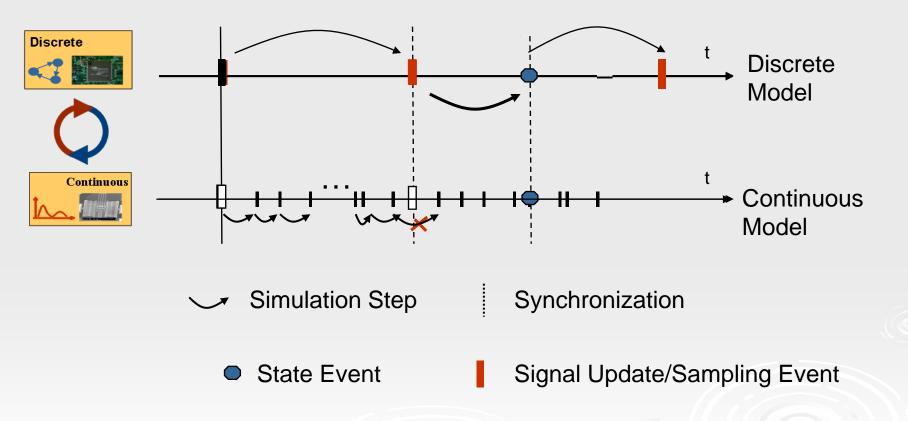
Concept Model	Time	Communication means	Processes activation rules
Discrete	Advances discretely	Set of events	Processes are
	(constant intervals)		sensitive to events
Continuous	It advances by	Piecewise-	Processes are
	integration steps (IS)	Continuous signals	executed at each IS

- Events exchanged between the continuous and the discrete models
 - State events
 - Sampling events
 - Update signal events

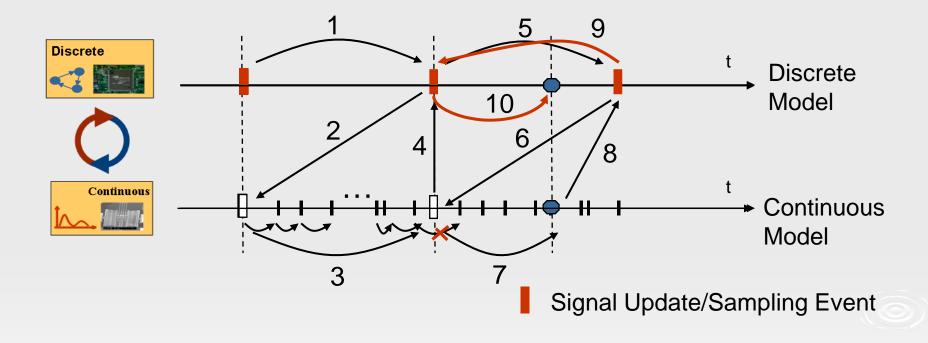


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Continuous/Discrete Systems - Global Simulation -



Rollback-Based Continuous/Discrete Synchronization Model

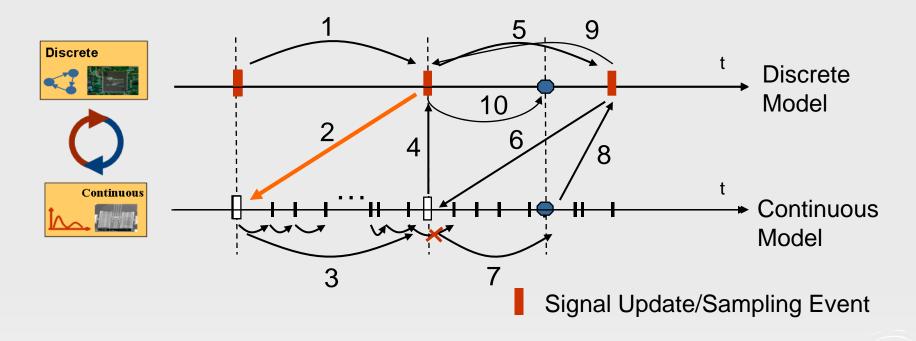


Synchronization Operational Semantics

- DEVS (Discrete Event Systems specification) and DESS (Differential Equation Systems Specification) formalisms
 - Set of rules Premise respecting the actions of Conclusion
 - Discrete model
 - Continuous model
 - Discrete model interface
 - Continuous model interface
- Timed Automata
 - Classical finite state automata with clock variables and logical formulas on the clocks (temporal constraints)

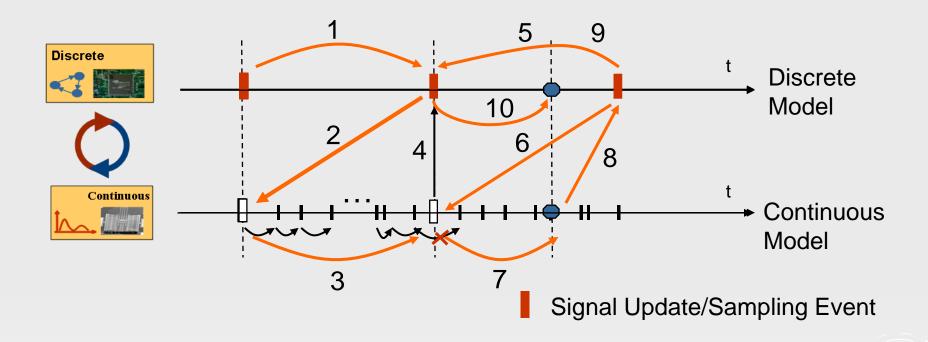


Synchronization Operational Semantics using DEVS



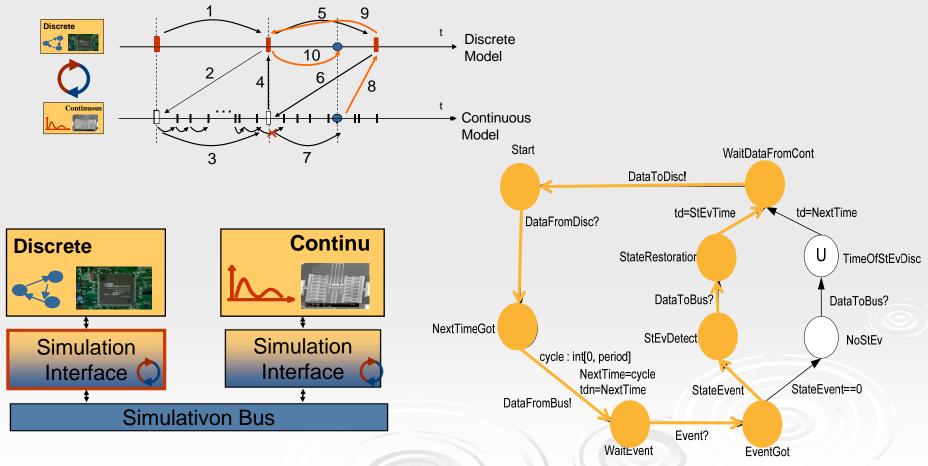
$$\frac{synch = 1 \land flag = 0 \land q = \delta_{ext}(q)}{q \xrightarrow{(DataFromBus, t_a(s_{dk}))?; synch := 0}} q$$

Synchronization Operational Semantics using DEVS

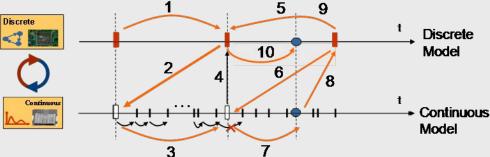


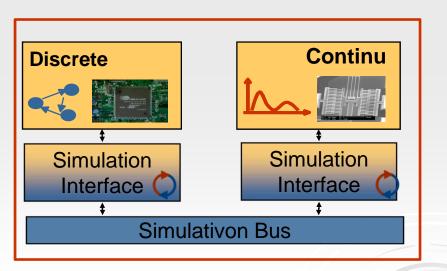
Set of rules for the overall synchronization model

Based on timed-automata – use of UPPAAL tool



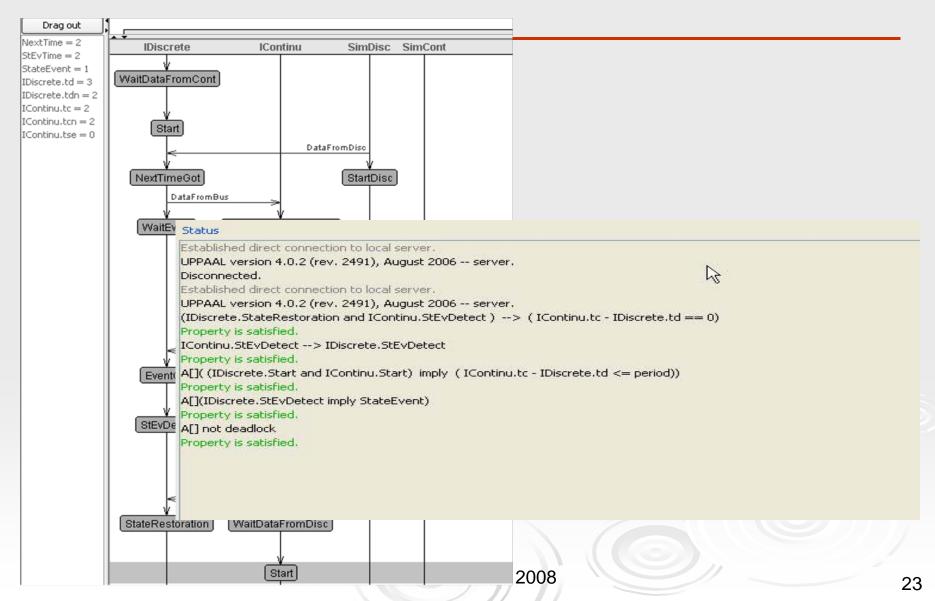
Based on timed-automata – use of UPPAAL tool





- Global Verification requires
 Timed Automata for
 - Discrete Simulator
 - Continuous Simulator
 - The Continuous
 Simulator Interface
 - The Discrete Simulator Interface

- Properties verified
 - Absence of deadlock
 - Timing synchronization
 - Detection of all state events
 - No false state events
 - Respect of causality principle



Conclusions

- Challenges for global validation of continuous/discrete systems
 - Definition of global execution models
 - Automatic generation of simulation interfaces
- Operational semantics definition and formal representation and verification
 - Key step for generalization and library definition

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